

Memory management

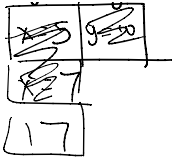
let f n =
 let l = [0; 1; 2; ...; n] in
 fun i → List.nth (l) i ← Closure
 let g = f 100
 gi = List.nth i
 g 3;; g 4;;

let f n =
 let l = [ref 0; ref 1; ...; ref n] in
 fun i → nth (l) i
 let g = f 10
 (g 3) (g 4) (g 8)
 lg = []
 g 3
 let h = f 10
 lh = []
 h 3
 (g 3) := 8

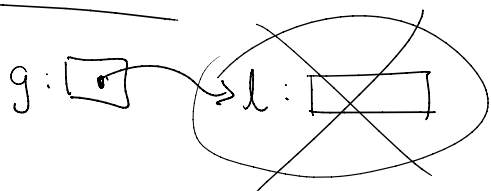
let x = 5 in
 ...
 let x = 5 in ←
 let y = x + 5 in ←
 let x = x + 2 in ←
 x + y

Question: What values are "escaping" this computation?

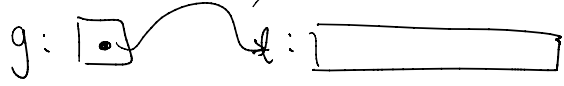




→ let g = f 10



→ let g = f 100



Stack / Heap

let (x) = 5 in

x + 3

- Variable does not escape
- So, it has a limited lifetime

→ Allocate it on the stack.

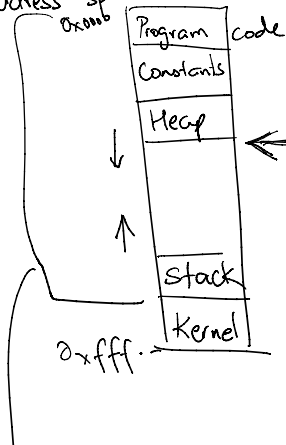
let g = f 100

- g escapes
- Its lifetime is not known "a priori"

→ Allocate on the heap.

Global Interpreter Lock.

Address space

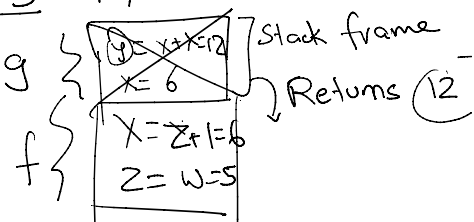


let g x = let y = x + x in y

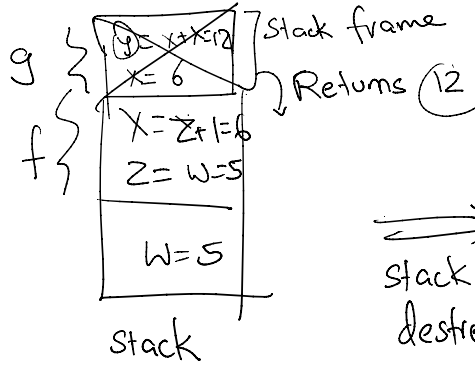
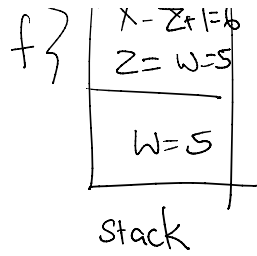
let f z = let x = (z + 1) in g(x) + g(x)

let w = 5 in

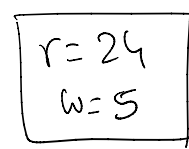
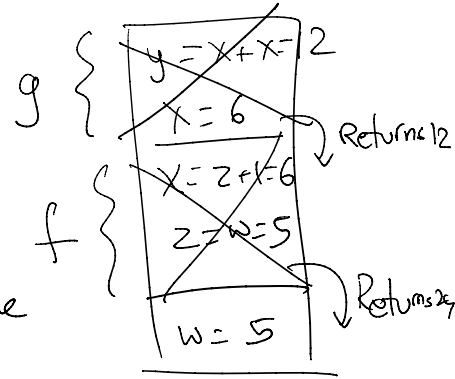
let r = f(w) in
w lives on



$0x111$
 w lives on
 128TB on x86-64
 Linux kernel.



Returns 12
 Returns 24
 stack frame destroyed.
 Memory reused



Memory management in C:

$\text{void} * \text{malloc}(\text{size}_t \text{ size})$ / oom

~~calloc~~
~~realloc~~
 Allocates "size" bytes & returns a pointer to them

$\text{free}(\text{void} * \text{ptr})$ ← Frees a pointer previously returned by malloc.

↓ ↑ Glibc / jemalloc / dlmalloc / ... / openBSD
 calls provided by the kernel

System calls provided by the kernel.

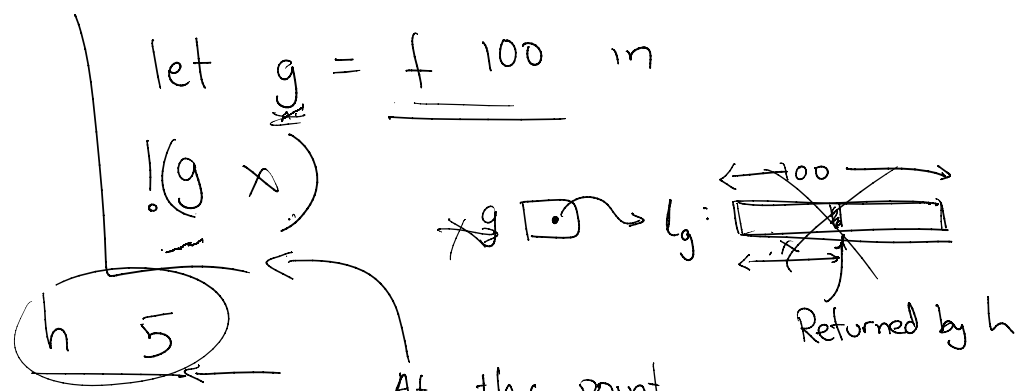
brk
sbrk

mmap
munmap

Garbage Collection

let f n =
 let l = [ref 1 ; ref 2 ; ... ; ref n] in
 fun i -> nth l i.

let h x =



At this point,
nobody can access
g any more.
Therefore, both g & lg can be freed.

(h in)

(h 10)

Therefore, both g & lg can be freed.



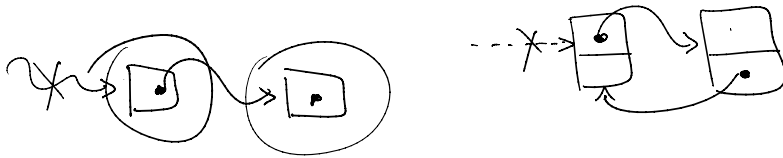
Question: Under what conditions can an object be freed?

Proposal 1: If nothing points to it.

Q1: Is it safe?

Repeat until exhaustion.

Q2: Is it complete?



Reference counting.

Proposal 2: Mark-&-Sweep Garbage Collector.



- Follow all pointers coming out of the stack.

- Do a DFS, mark all heap locations that you visit.

- Make a pass over the heap deallocate all cells which

have not been marked.

GC-pause Garbage collection pause	GC-pressure.
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Rust has no GC.

Pauseless GC. vs. stop-the-world.

Epsilon GC in new JVMs.