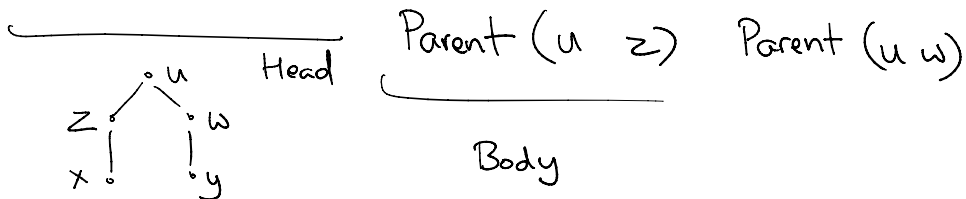


## Datalog

Cousins (x y) :-  $\overbrace{\text{Parent}(z x) \text{ Parent}(w y)}$



Path (x z) :- Path(x y) Edge(y z)

Path (x y) :- Edge(x y).

"All-pairs reachability"

Transitive  $\left\{ \begin{array}{l} a \leq b \quad b \leq c \\ \hline a \leq c \end{array} \right.$

Transitive closure : Begin with a seed relation.

"Make it transitive"

- Look for violations of the transitivity rule.

- Whenever  $\exists x y z$  s.t.

$x \leq y$  &  $y \leq z$

If it is not yet the case  
that  $x \leq z$ ,

then just "throw it in".

Add  $x \leq z$ .

Parent ( $x$ $y$ )	}	Ancestor ( $x$ $y$ )
$x$ is a parent of $y$ .		$x$ is an ancestor
$x \geq y$		of $y$ .

---

The Parent relation is not  
transitive

"My grandparent is not  
my parent."

---

The Ancestor  
relation is transitive.

— My grandparent is  
an ancestor of mine.  
— His grandparent is  
also an ancestor of  
mine.

---

Equivalently, the transitive closure of a  
relation  $R$  is the smallest relation  $R^*$

- st.
- ①  $R \subseteq R^*$
  - ②  $R^*$  is transitive.

---

Plan

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## Plan

- ① Become comfortable with Datalog  
"Writing queries".
- ② Demystify the language  
Evaluation algorithms
- ③ "The Problem of Negation".

Stratified programs

Aggregation ↙

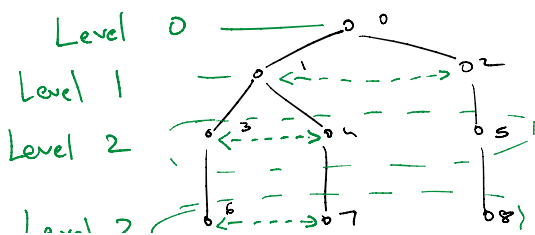
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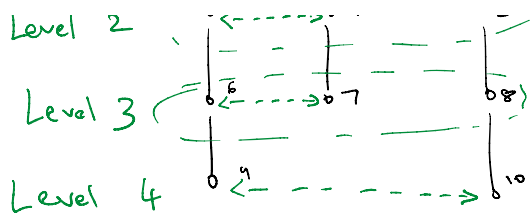
Write a program to find all pairs  
of people who are in the same  
generation of the family.

---

Parent (x y) = "x is a parent of y"

Samegen (x y) = "x & y are in the same  
generation"





Rule 1: "We are in the same generation

if our parents are in the same gen."

↓

samegen(x y) :- samegen(u v),  
 parent(u x), parent(v y).

Rule 2: Base case.

samegen(x x) :- person(x).

↑

We just invented this relation.

person(x) :- parent(x, \_)

person(x) :- parent(\_, x).

① samegen(x, y) :- samegen(u, v), parent(u, x), parent(v, y).

samegen(x, y) :- parent(z, x), parent(z, y), x != y.

← Base case of

samegen(x, y) :- parent(z, x), parent(z, y), x != y.

Alternative program which avoids

Base case of the induction.

samegen(x x)

Assume not. Say we have a tuple  
samegen(x x).

Ask the prosecutor: "Show your evidence!"

— Could the last step in the stack trace be Rule 2?

No, because  $x \neq y$ .

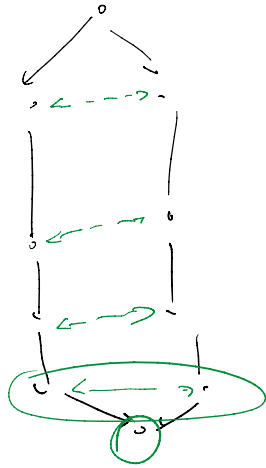
— Could the last step in the stack trace be Rule 1?

samegen(x y) :- samegen(u v)

parent(u x) parent(v y),  $x \neq y$ .

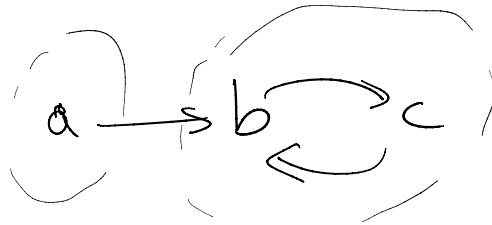
samegen(x x)  $\Leftarrow$  samegen(u v)

parent(u, x)    parent(v, x)

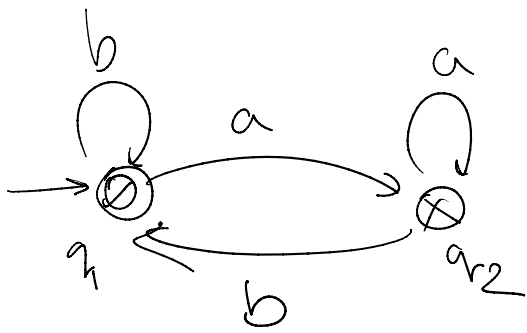


Q2: Find if two nodes in a graph are in the same SCC.

Edge	
a	b
b	c
c	b



$scc(u, v)$  if there is a path from  $u$  to  $v$  & a path from  $v$  to  $u$ .



Write a <sup>Datalog</sup> program which executes this DFA.

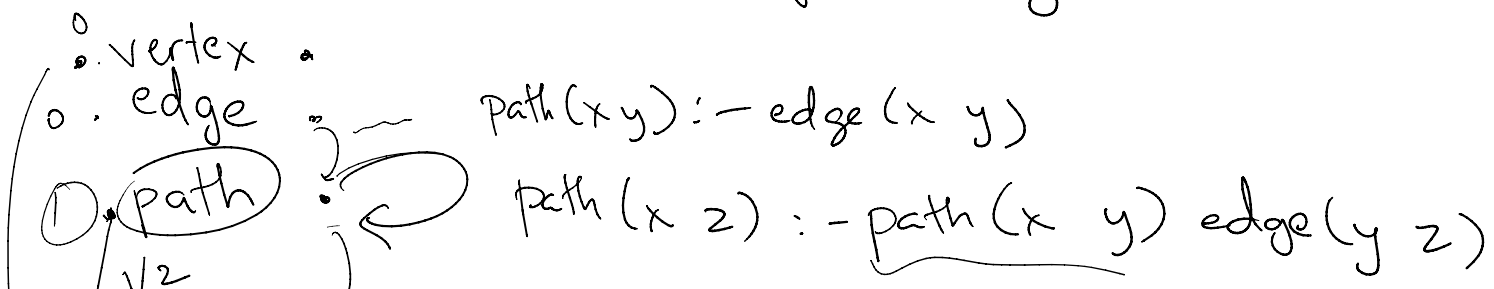
Index	0	1	2	3	4	5	6	<del>7</del>
Char	a	a	b	b	a	a	b	<del></del>

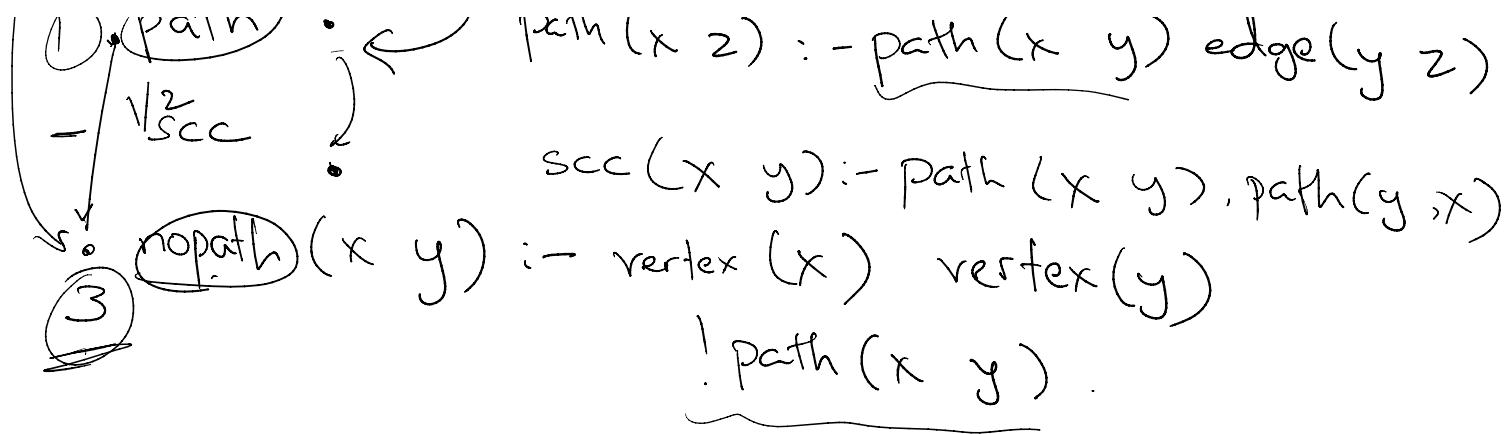
Negation assumes "closed world".

The closed world assumption states:

"All that is true is known to be true."

Souffle employs "stratified negation"





Paradox because ! operator in a cycle.

Souffle tries to check if application of ! forms a DAG.