

## Structuring Data with Types

Basic types 1 : int, float, char, bool

Basic types 2 : unit, function ( $\_\rightarrow\_\_$ )

Standard types 1: string, pair, tuple,  
list, exn, option

Standard types 2: array

Custom types: variants, records  
↓  
Algebraic Data Types (struct in C)

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type	status	=	Valid		Invalid
			↑		
			"or"		

- Values of type status can be either Valid or Invalid
- "That is all they can be."

- That is all they can be.  
 "Closed world assumption".
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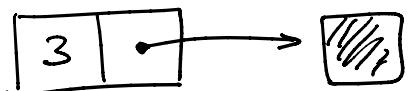
## Recursive Types

type ourlist = EmptyList

|  
Not Empty List of int \* ourlist.  
NEL



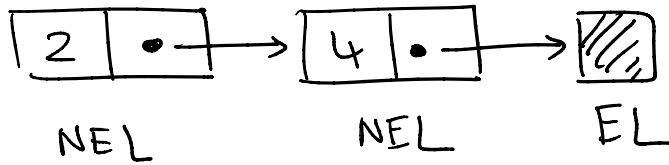
EmptyList



NEL



EL



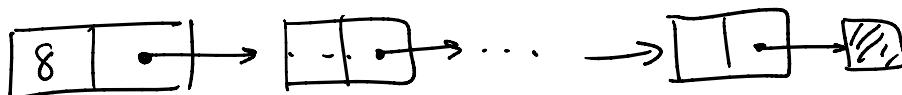
NEL

NEL

EL

ourlist

:




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let cons hd tl = NEL (hd, tl)

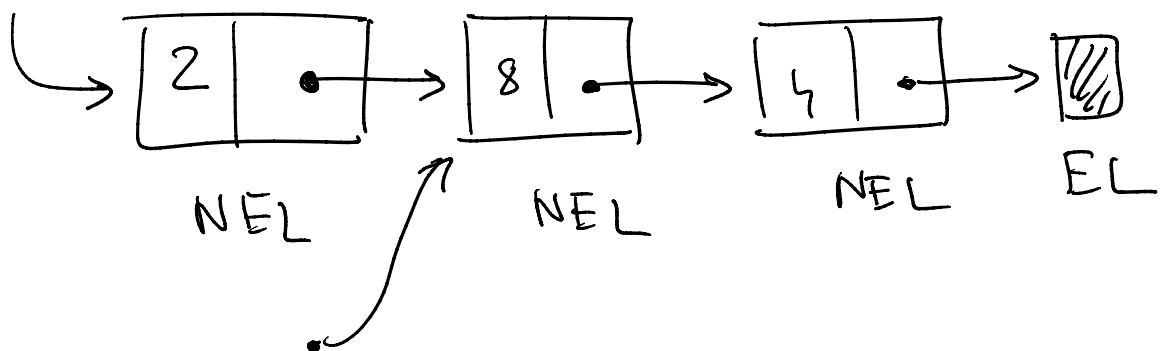
let car l = match l with

```

| EL → ???
| NEL(hd, _) → hd
let cdr l = match l with
| EL → ???
| NEL( _, tl) → tl

```

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Computing the length of a list

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```

let ourlen l =
match l with
| EL → 0
| NEL( _, tl) → 1 + ourlen tl

```

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```

let rec ourlen l =
match l with
| EL → 0
| NEL( _, tl) → 1 + ourlen tl

```

$| \text{NEL}(-, \text{tl}) \rightarrow | + \text{ourlen } \text{tl}$

???

Attempts to  
resolve ourlen in  
a scope before the  
present fn is defined

$| \text{NEL}(-, \text{tl}) \rightarrow | + \text{ourlen } \text{tl}$

Adds ourlen to  
scope before  
present definition

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To be able to print ADTs

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Magic: Use  $[@@ \text{ deriving show}]$   
after definition

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Concatenating lists

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let rec concat  $l_1$   $l_2$  =

match  $l_1$  with

$| [] \rightarrow \underline{l_2}$

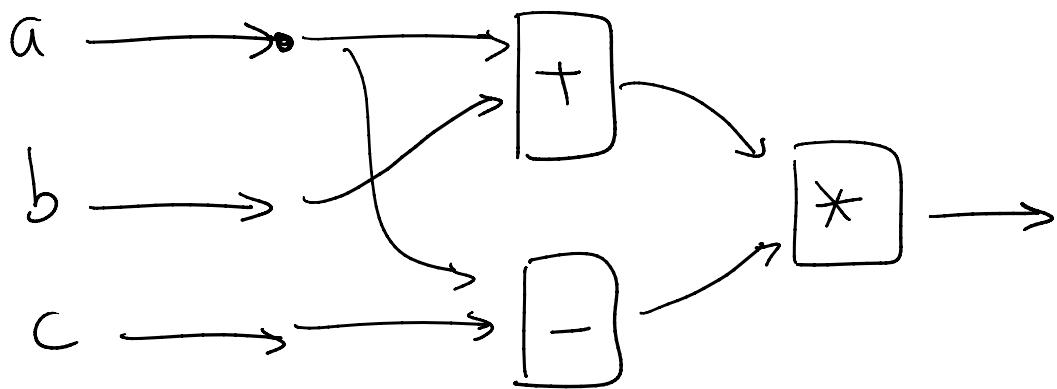
$[2; \underline{9}; \underline{3}] [4; \underline{2}; \underline{8}]$   
 $\Rightarrow [2; \underline{9}; \underline{3}; 4; \underline{2}; \underline{8}]$

$$| [] \rightarrow \underline{l_2} \Rightarrow [2; \underline{9}; \underline{3}; \underline{4}; \underline{2}; \underline{8}]$$

$$| \text{hd} :: \text{tl}_1 \rightarrow \underline{\text{hd} :: (\text{concat } \underline{\text{tl}_1}, l_2)}$$

Recursive definition terminates  
because first argument is  
strictly decreasing.

$$\text{let funny-foo } a \text{ } bc = (a+b) * (a-c)$$



- No matter how big  $a, b, c$  are, this circuit will always be 3 gates big.

- In the absence of recursion, all computations result in circuits of bounded size.

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let sum n =

$$\sum_{i=0}^n i$$

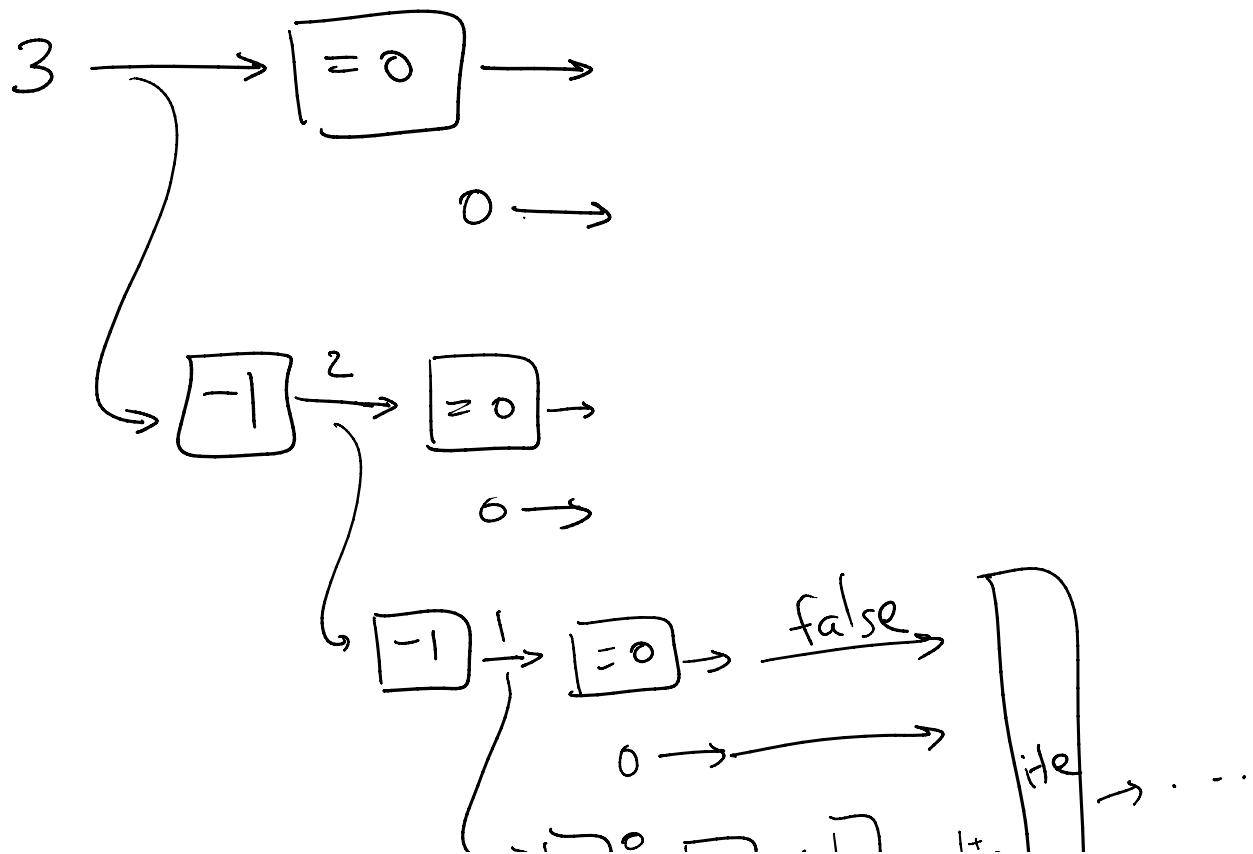
if  $n = 0$  then 0

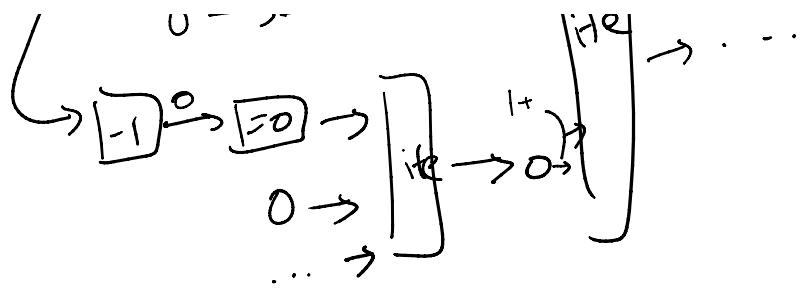
$$0 + 1 + 2 + \dots + n$$

else  $n + \text{sum}(n-1)$

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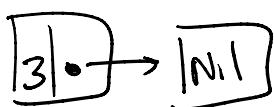
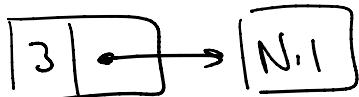
sum 3





Recursion provides the ability to build circuits that are arbitrarily large.

Structural equality	vs.	Physical equality
$=$ , $\langle \rangle$		$==$ , $!=$
$\boxed{3} \rightarrow \boxed{\text{Nil}}$		(Don't worry about these for now.)



- ① Checks if values have the same shape
- ② If so, recurses
- ③ Might not terminate for cyclic structures

## Merging sorted lists

```
let rec merge l1 l2 =
(* Assume that l1 is sorted *)
(* Assume that l2 is sorted *)
(* We want to merge them into a single
sorted list *)
match l1, l2 with
| [], _ -> l2
| _, [] -> l1
| hd1 :: tl1, hd2 :: tl2 ->
if hd1 < hd2
then hd1 :: merge tl1 l2
else hd2 :: merge l1 tl2
```

## Quick-and-dirty QuickSort

```
let rec qsort l =
match l with
| [] -> []
| hd :: tl ->
let tl1 = List.filter ((>) hd) tl in
let tl2 = List.filter ((<=) hd) tl in
let stl1 = qsort tl1 in
let stl2 = qsort tl2 in
stl1 @ [hd] @ stl2
```

Sorted versions  
of  $tl_1$  &  
 $tl_2$  resp.

All elements of  $tl$  which  
are smaller than  $hd$

All elements of  $tl$  which  
are at least as large as  $hd$

$tl_1$  &  $tl_2$  contain  
strictly fewer elements  
than  $l$ .

Concatenate everything

