

Memory Management

Program

Language runtime (Pointers, memory allocation, garbage collection)

Operating system (Virtual memory, page tables, swapping, ...)

Hardware

(Memory controller, Direct Memory Access (DMA)

Translation Lookaside Buffer (TLB),
...)

```
let f n =  
  let l = Base.List.range 0 n in  
  fun i -> List.nth l i
```

This call to f allocates
a list l

```
-( 16:13:20 ) < command 5 > { counter: 0 }-  
utop # let g = f 10;  
val g : int -> int = fun -  
-( 16:13:44 ) < command 6 > { counter: 0 }-  
utop # g 3;  
- : int = 3  
-( 16:13:57 ) < command 7 > { counter: 0 }-  
utop # g 5;  
- : int = 5  
-( 16:14:13 ) < command 8 > { counter: 0 }-  
utop # g 7;  
- : int = 7
```

All calls Refer to the same l.

```
let g2 = f 100
```

```
g2 73
```

Creating g_2 does not
invalidate g

There must be two independent
l-s in memory.

Jobs of the language runtime

- Memory management

- What things are in memory?

- What memory is used?

- What memory is available?

```
let rl = ref (Base.List.range 0 10)
```

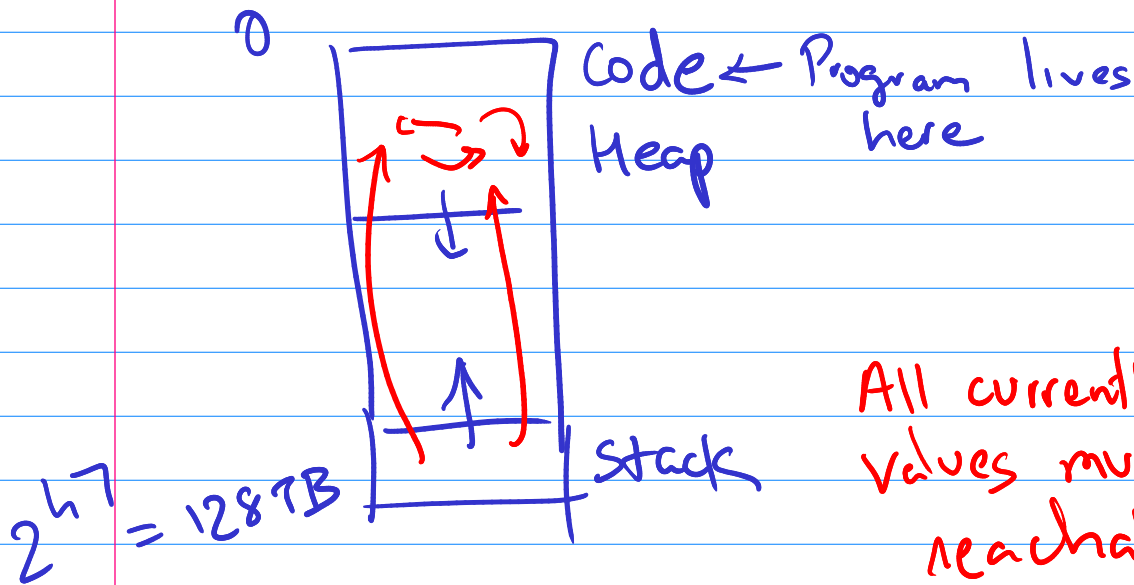
```
rl := (List.map Int.succ (Base.List.range 0 200))
```

Old list not reachable any more.

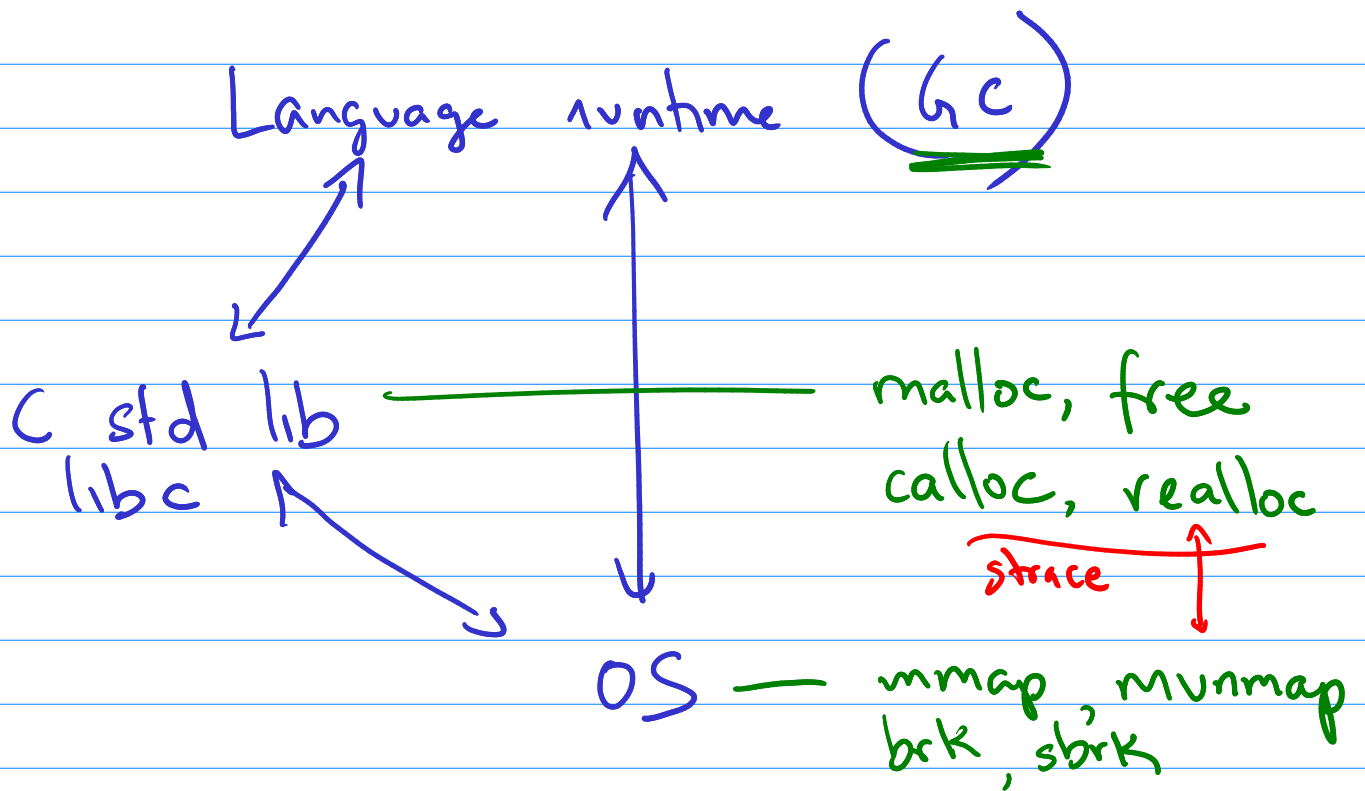
Memory Can be reclaimed!

Virtual address space

2^{64}



All currently accessible values must be reachable from the stack.



Memory management / Garbage collection

- ① Where to allocate an object?
- ⇒ ② When is it safe to deallocate an object?
- ③ When to actually deallocate?

Proposal: An object is safe to deallocate
if nothing points to it.

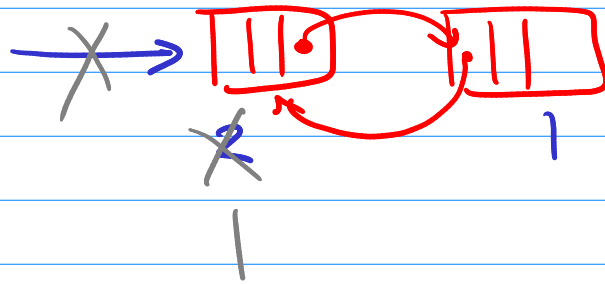
(Ref-counting) → Python

→ C++ auto ptrs

+ simple to implement (careful about concurrency!)

+ fast

- but it can leak memory!



"problem of circular references:"

Proposal 2 : Mark-&-sweep garbage collector

"Pauseless GCs"

Step 0 : Pause the program

"Stop-the-world GCs"

Phase 1 : Start from the "root set"

- Follow the chain of pointers

- Mark all visited objects as "live"

Phase 2 : Go over all currently allocated objects

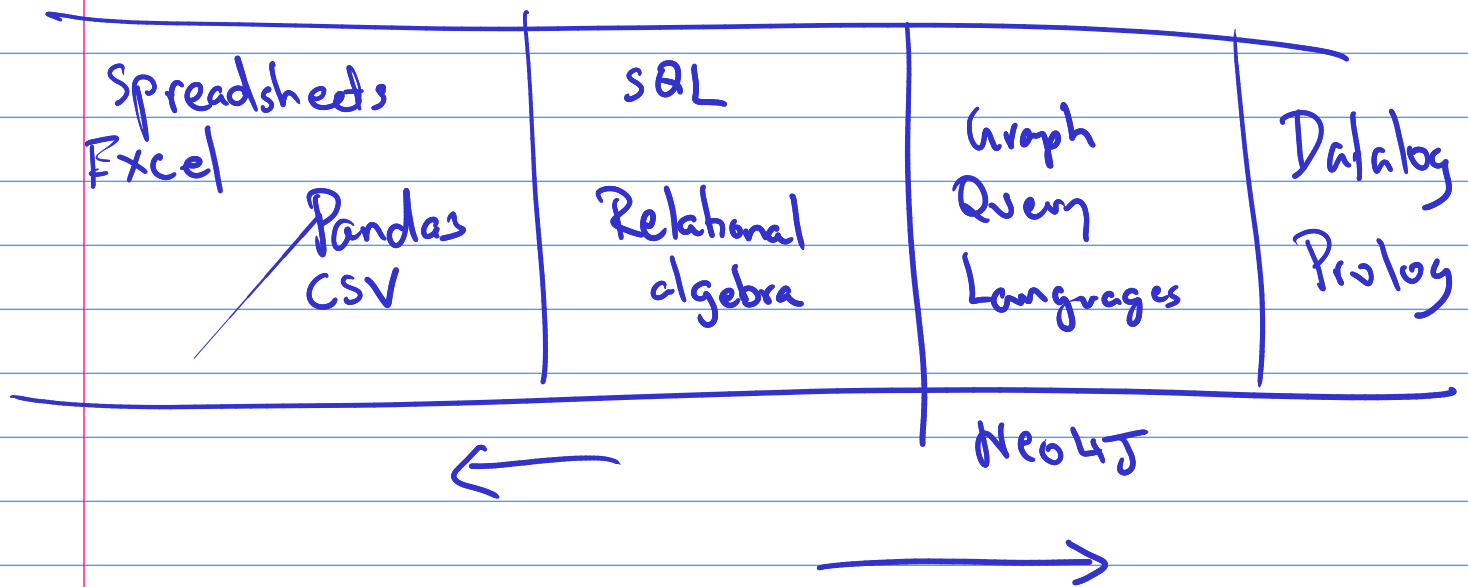
- If not live, then mark them

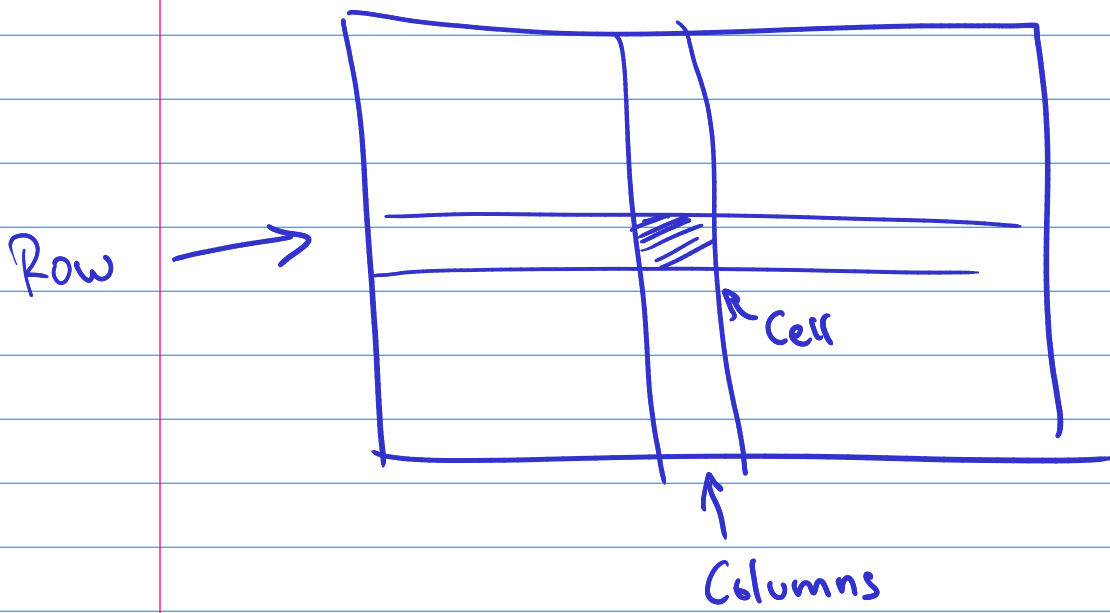
as candidates for deallocation.

Generational garbage
collection

Newly created objects more
likely to become unreachable
than old objects.

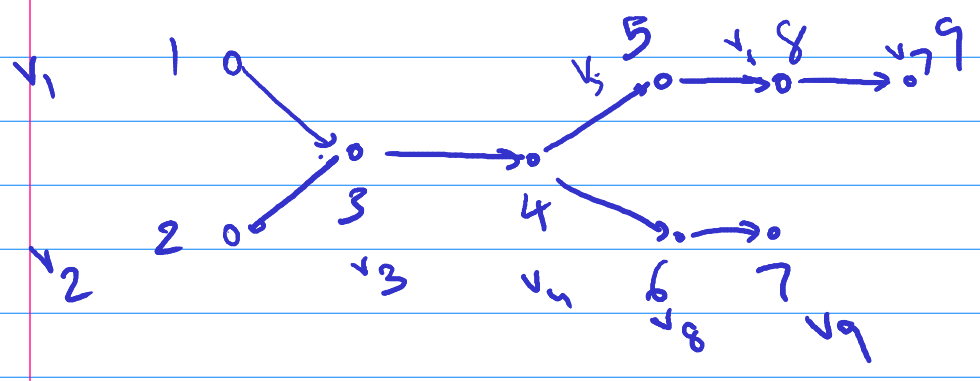
Unit 3 : Programming With Relations





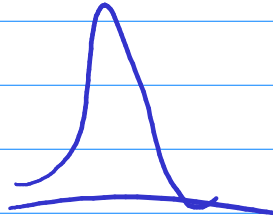
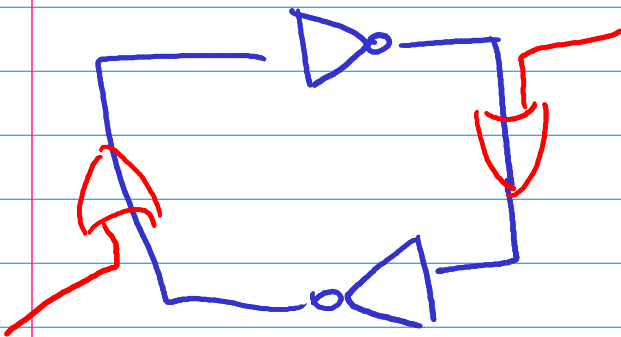
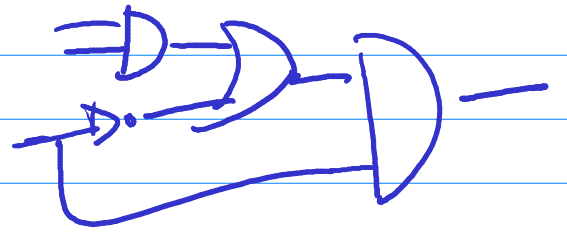
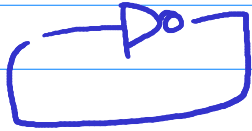
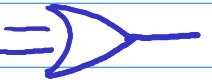
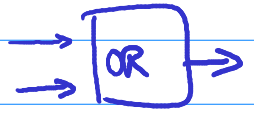
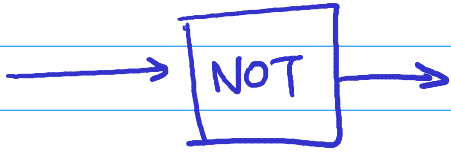
Topological Sort

- Check if a graph is a DAG (Directed Acyclic Graph)
- Emit a proof that it is a DAG



- | |
|-----------------|
| $0 < v_1$ |
| $0 < v_2 \dots$ |
| $v_1 < v_3$ |
| $v_2 < v_3$ |
| $v_3 < v_4$ |
| $v_4 < v_5$ |
| $v_4 < v_6$ |
| $v_5 < v_8$ |
| $v_6 < v_7$ |
| $v_8 < v_9$ |
| \vdots |

$$\min(v_1 + v_2 + \dots + v_9)$$



DAGs → Spreadsheets /
Makefiles / Build systems
Combinatorial circuits

Q1: Is there / what is the equivalent of
pointers?

Q2: What is the counterpart of a dictionary?

Q3: Can we do simple "real" computations
in Excel.