

① Paper presentations on Feb 8

Model counting, ML-enabled decision heuristics

QBF solvers

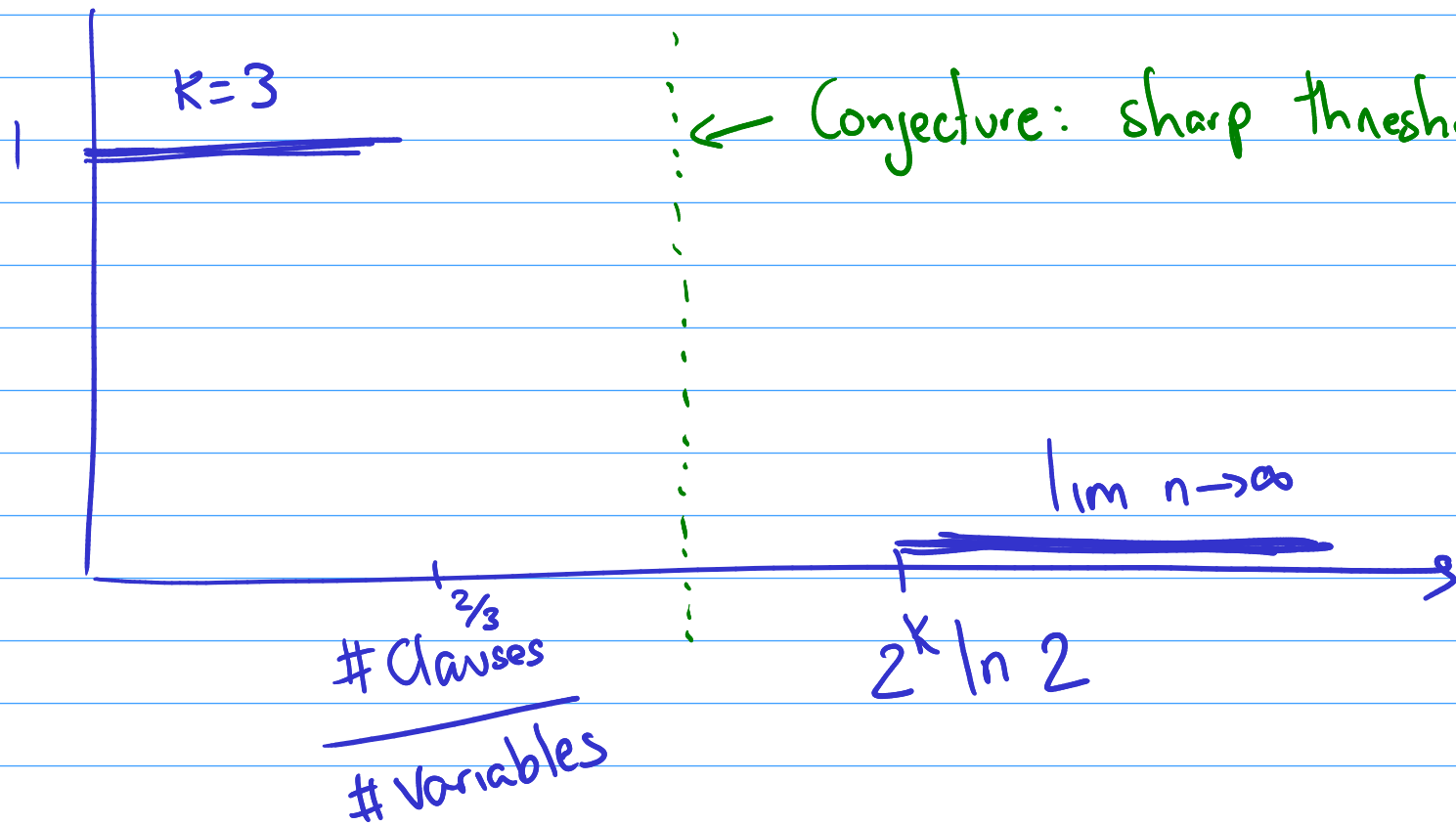
## WalkSAT

- ① Guess an assignment
  - ② Repeat until assignment  $\models \varphi$  :
    - Pick an unsatisfied clause, uniformly at random
    - Pick a literal, uniformly at random
    - Flip it.
- $l_1 \vee l_2 \vee \dots \vee l_k$
- "Models"  
↙

## Random - CNF

- ① Fix  $n$  variables, clauses of length  $K$ .
  - ② Decide that formula has  $m$  clauses:  
Sample clause (with replacement).
- $\binom{n}{K} 2^K$   
possible clauses
-

Probability  
of satisfying  
assignment



## Erdos-Renyi graphs $G(n, p)$

For each pair  $(u, v)$  of vertices (undirected)  
with prob.  $p$ , insert edge  $(u, v)$ .

Disconnected almost surely, if  $p < \frac{1}{n}$

Many vertices in same SCC, if  $p > \frac{1}{n}$ .

# Satisfiability Modulo Theory

int  $x_1$   $x_2$   $x_3$  := input();

if ( $x_1 \neq x_2$ )

if ( $x_2 = x_3$ )

if ( $x_3 = x_1$ )

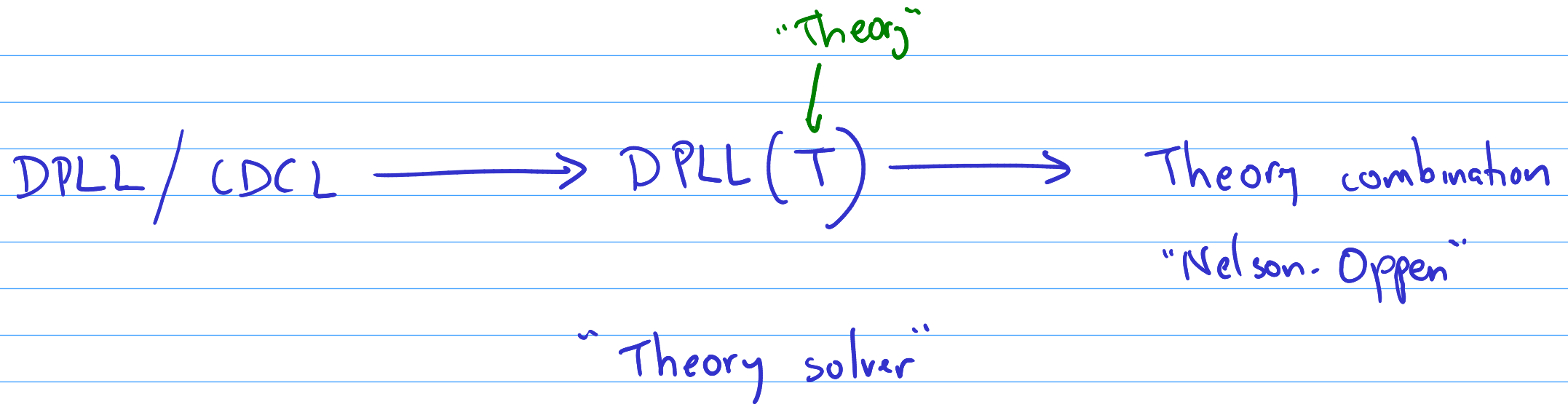
crash!

$\exists x_1 x_2 x_3$  st.  $(x_1 \neq x_2) \& x_2 = x_3 \&$   
 $x_3 = x_1?$

int i = j

arr[i] := 10;

if (arr[j] < 8) crash!

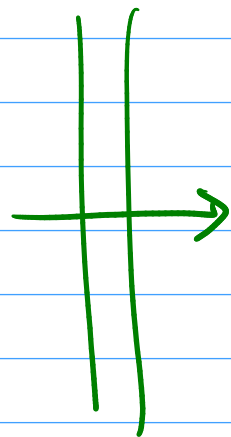




## A solver for Linear Real Arithmetic

Q: Do there exist real numbers  $x$  &  $y$  s.t

$$\begin{aligned} & x + 2y \leq 5 \\ & \& x - 8y > 3 \\ & \& \neg (x + 4y < 8) \end{aligned}$$



$$\begin{aligned} & x + 2y \leq 5 \\ & x - 8y > 3 \\ & x + 4y \geq 8 \end{aligned}$$

Submit to favorite LP solver

$\exists x \cdot \exists y \cdot$  \_\_\_\_\_

$\forall x \cdot \exists y \cdot \forall z \cdot \exists w$  \_\_\_\_\_

Eliminate  $w$   
innermost variable  
"Quantifier elimination"

Conjunction of  
linear equations & inequalities

$$x + 8y \leq 3, \quad x - 2z = 9, \quad \dots$$

$$x^2 - 4yz \geq 3$$

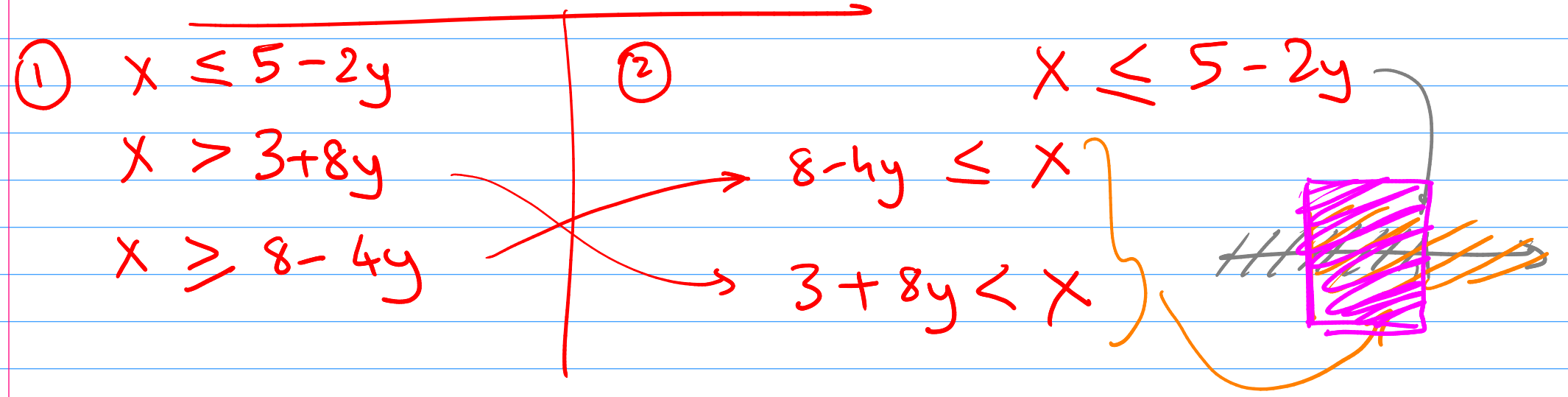
$$x = y$$

$$x \geq y$$

$$y \geq x$$

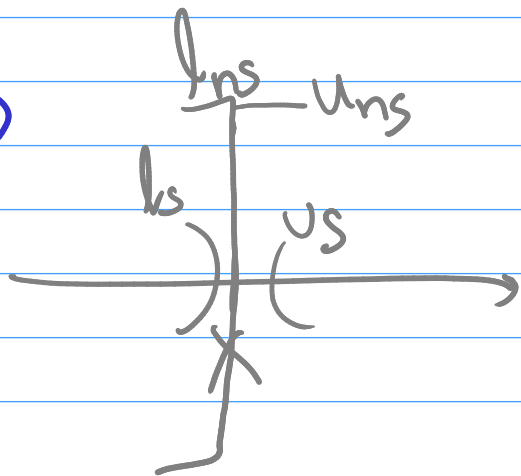
Don't worry about equalities

Q  $\exists y. \exists x. x + 2y \leq 5$   
 $x - 8y > 3$   
 $x + 4y \geq 8$



$$\begin{array}{ccc}
 \exists x \cdot & x < u_s & \\
 & x \leq u_{ns} & \\
 l_{ns} \leq x & & \\
 l_s < x & & 
 \end{array}
 \longleftrightarrow
 \begin{array}{ccc}
 l_{ns} < u_s & & \\
 l_{ns} \leq u_{ns} & & \\
 l_s < u_s & & \\
 l_s < u_{ns} & & 
 \end{array}$$

Proof  $\Rightarrow$



$\Leftarrow$  If there is a gap, then you can pick  $x$

$$\forall x \text{ ————— } \Leftrightarrow \neg \cdot \underbrace{\exists x \cdot \neg \text{ —————}}_{\text{—————}}$$

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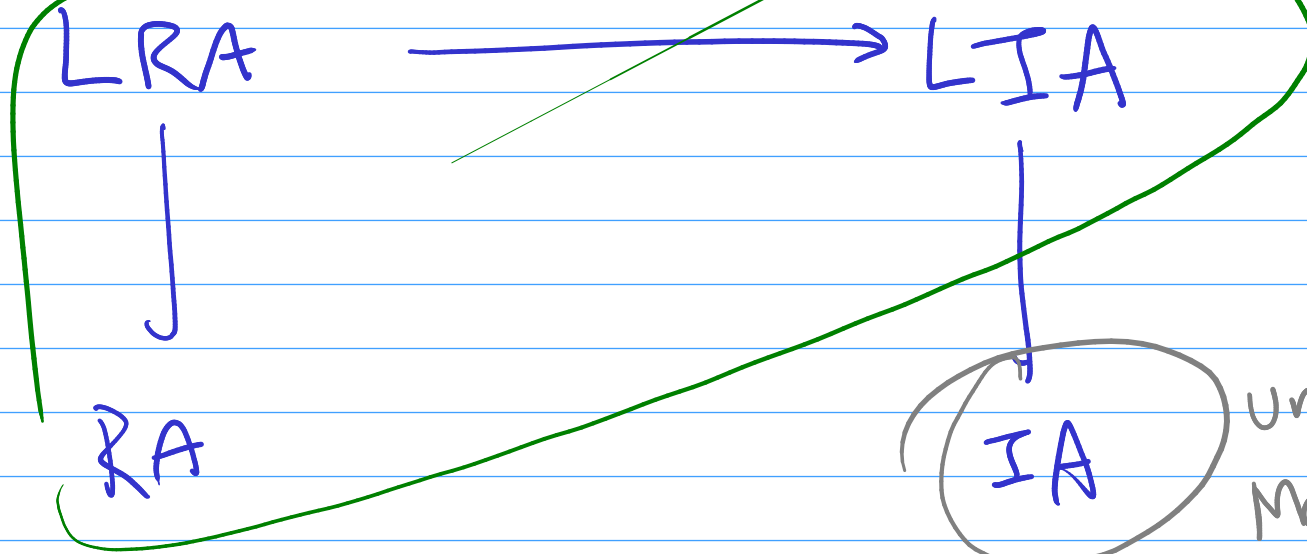
$$\forall x \cdot x < 3y \ \& \ y < 8x$$

$$\neg \cdot \exists x \cdot \neg (x < 3y \ \& \ y < 8x)$$

$$\neg \exists x \cdot (x \geq 3y \ \text{or} \ y \geq 8x)$$

FO

Decidable / Quantifier elimination



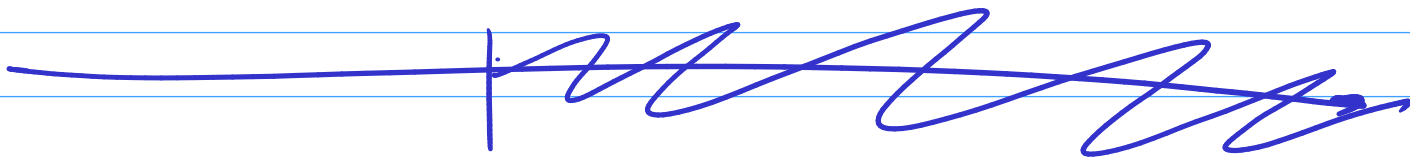
undecidable  
Matiyasevich's theorem  
Hilbert's 10<sup>th</sup> problem

$$\forall x. \exists y. x^2 = y : \text{sat}$$

$$\forall_{P \Rightarrow X} x. \exists_{P \Rightarrow y} y. x = y^2 : \text{Unsat}$$

First order  $\forall x. \exists y. x < y$  ? Yes :  $y = x + 1$

Second order  $\forall X. \exists y. \forall x \in X. x < y$  ? No



There exist exactly two values of  $x$  s.t.  $\varphi(x)$

$$\exists x. \exists y. \varphi(x) \& \varphi(y) \& \forall z. \varphi(z) \Rightarrow (x=z \text{ or } y=z) \\ \& x \neq y$$



Ex:

$\exists x \in \mathbb{R} \quad 3 < x \quad \text{and} \quad x < 4 \quad ? \quad \text{Yes}$

$\exists x \in \mathbb{Z} \quad 3 < x \quad \text{and} \quad x < 4 \quad ? \quad \text{No}$

$\exists x. \exists y. z \quad x = y + z \quad \text{and} \quad 3 < x \quad \text{and} \quad y + z > 3 ?$

$$\underline{\exists x. y = x + x + x}$$

$$\varphi(y)$$

$$y \bmod 3 = 0$$

Theorem 6.12

from Sipser's

text book