

$$u - 8 \leq 16 + w \iff u - w \leq 24$$

Difference Logic

$$\begin{array}{ccc} & \nearrow & \nwarrow \\ u - 8 \leq v & & v \leq 16 + w \\ \uparrow & & \uparrow \end{array}$$

Ex: $\exists u, v, w$ s.t. $u - v \leq 8$ $v - w \leq 16$ $u - w \geq 32$?

$$\Rightarrow u - w \leq 24$$

$$u - w = 0 \iff u - w \leq 0 \text{ and } u - w \geq 0$$

Difference logic : ① All variables range over \mathbb{Z} (or \mathbb{R})

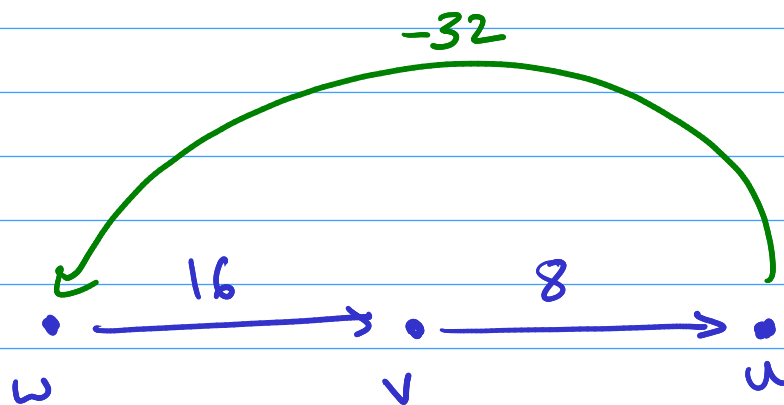
② Formula is a conjunction of constraints

③ Each constraint is of the form $u - v \leq c$,
for some variables u, v
constant c .

$\exists u, v, w$ s.t. $u - v \leq 8$ $v - w \leq 16$ $u - w \geq 32$?

$$-(u - w) \leq -32$$

$$w - u \leq -32$$



① Vertices = all variables

② $u \xrightarrow{c} v$ edge whenever constraint $v - u \leq c$

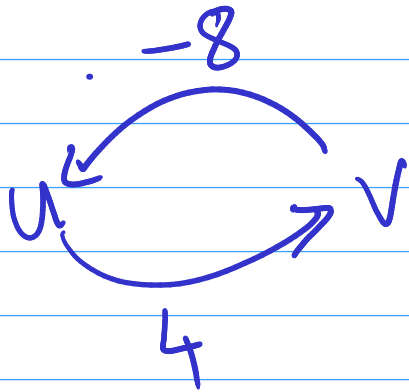
From this graph, we can recover the original formula.

\therefore The graph itself encodes satisfiability "in some form".

Conjecture Difference logic formula φ is satisfiable

iff no negative wt. cycles in G_φ .

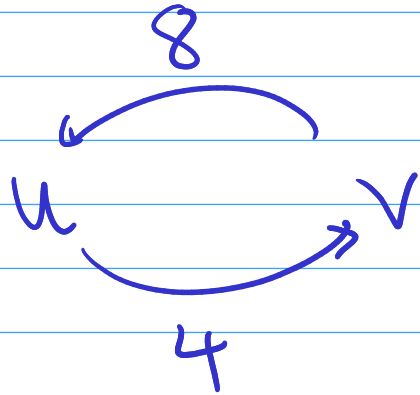
Ex 1:



$$\left. \begin{array}{l} v - u \leq 4 \\ u - v \leq -8 \end{array} \right] \Rightarrow \underline{\underline{0 \leq -4}}$$

Ridiculous

Ex2:

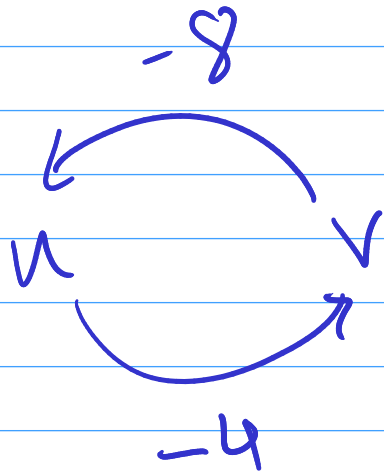


$$v - u \leq 4$$

$$u - v \leq 8$$

Satisfiable, set both to 0

Ex3:



$$v - u \leq -4$$

$$u - v \leq -8$$

$$0 \leq -12$$

Unset

Ex 4 : $u \xrightarrow{18} v \xrightarrow{27} w$

$$v - u \leq 18 \quad w - v \leq 27$$

$$w - u \leq 45$$

Claim : If G_φ contains a path $u \overset{c}{\rightsquigarrow} v$

then $\varphi \Rightarrow v - u \leq c$

Claim: If G_φ contains a negative weight cycle,

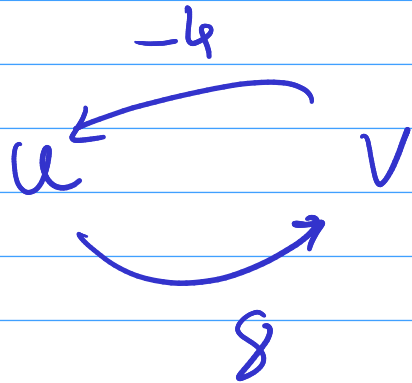
$u \overset{c}{\rightsquigarrow} u$, with $c <_+ 0$,

then $\varphi \Rightarrow u - u \leq c$

$\Rightarrow 0 \leq c$ Preposterous

& φ is unsat

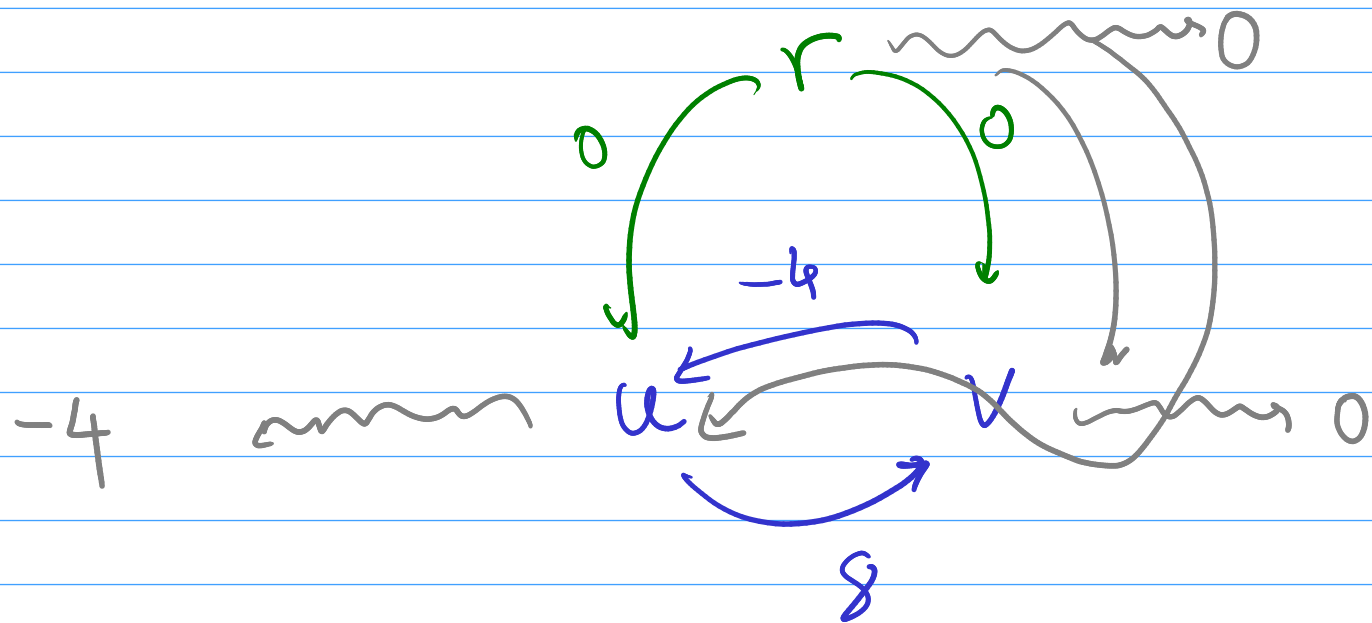
Ex 4:



u		-4	-3	-2	\dots	0
v		0	1	2		4

$$v - u \leq 8$$

$$u - v \leq -4$$



Claim: Construct the assignment $m \{ u \mapsto \text{dist}(r, u) \mid u \}$
 $m \neq \emptyset$.

Theory of Equality

Ex: $\exists x y z \in \mathbb{Z} \cdot x=y \quad x \neq z \quad y=z?$ ^{Unset}

$\exists x y z \in \mathbb{R} \cdot x=y \quad x \neq z \quad y=z?$

$\exists x y z$ strings

$\exists x y z$ objects



Theory of equality

$\exists v_1, v_2, \dots, u_1, u_2, \dots \in \text{set of objects of your choosing}$
 $\varphi = v_1 = v_2 \text{ and } v_3 = v_4 \text{ and } \dots \text{ and } v_k = v_{k+1}$

and $u_1 \neq u_2$ and $u_3 \neq u_4$ and \dots and $u_m \neq u_{m+1}$

?

$\varphi = \exists x y z \in \mathbb{D} \cdot x=y \quad x \neq z \quad y=z?$



$G_\varphi = \{ u-v \mid \text{whenever } (u=v) \in \varphi \}$

Claim: If $u \neq v \in \varphi$ & u & v are in the same SCC in G_φ , then φ is unsat.

Equality with Uninterpreted Functions

Ex

$\exists f: \mathbb{Z} \rightarrow \mathbb{Z} \exists x \in \mathbb{Z}$

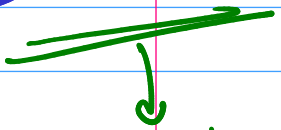
$f(x) = x$

and $f(f(x)) \neq f(x)$?

$f(x) \neq f(x)$

$f(x) \neq x$

$\exists f: \mathbb{R} \rightarrow \mathbb{R} \cdot \exists x \in \mathbb{R}$



'uninterpreted'

$f(x) = x$ and $f(f(x)) \neq f(x)?$

unsat

$\exists f: \text{string} \rightarrow \text{string} \cdot \exists x \in \text{string} \cdot f(x) = x$ and $f(f(x)) \neq f(x)?$

float $x_0 = \text{input}$

$x_1 = \text{sqrt}(x_0)$

$x_2 = \text{sqrt}(x_1)$

If $(x_1 = x_0 \ \& \ x_2 \neq x_1)$ then crash!