

Theories

- LRA LIA RA IA
- Difference logic
- Theory of equality / EUF

$$\exists x. x + 2.5 \leq 3.8 \quad \text{LRA}$$

$$\exists f. \exists x. y. f(x) = f(y) \text{ and } x \neq y \quad \text{EUF}$$

$$\exists x. y. x - y \leq 4 \text{ and } y - x \leq 8 \quad \text{IDL}$$

"non-logical"
 Set of symbols Σ } "Signature"
 Each symbol has an arity } ~~"Logic"~~ + variables
 } + "=" "logical"
 } + "∃", "∀"
 } + "Sentences" + "∨, ∧, ¬"

$$2 \cdot 0 + 4 \cdot 8 = 6 \cdot 8$$

$$3 \cdot 3 - 2 \cdot 6 = 1 \cdot 9$$

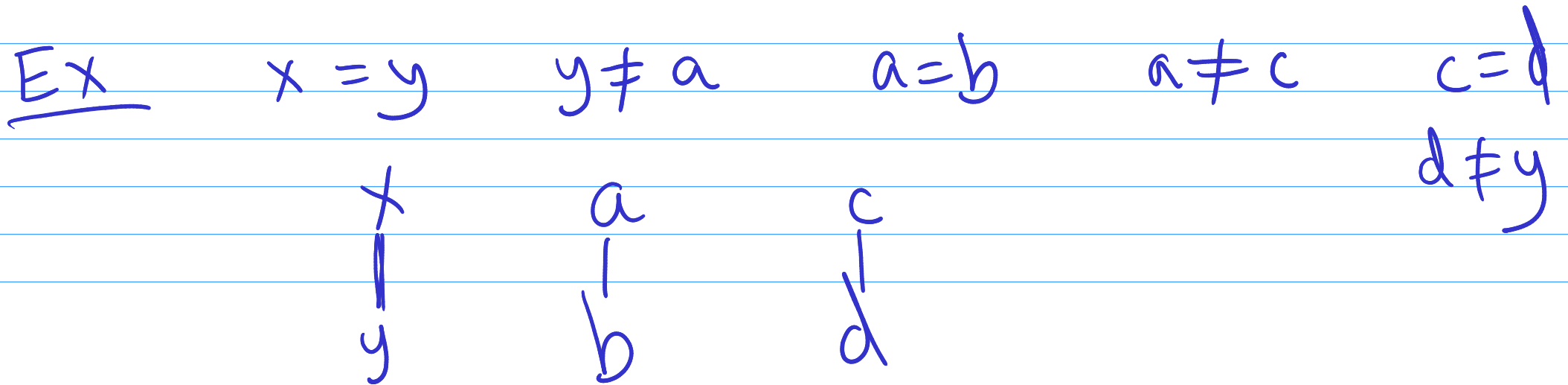
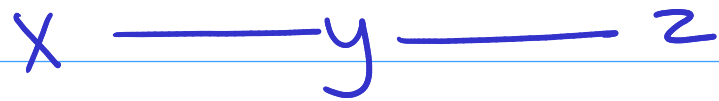
$$3 \cdot 3 \neq 4 \cdot 8$$

"Theory" is a set of sentences

$\exists x. \quad 3 < x \quad \text{and} \quad x < 4$ LIA

LRA

Ex : $x = y$ $y = z$ $x \neq z$



Equality with uninterpreted fns

Ex: $\exists f \exists x \cdot f(x) = x$ and $f(f(x)) \neq x$?

$x \xrightarrow{f} f(x) \rightsquigarrow f(f(x))$

Ex $\exists f \cdot \exists x y \cdot x=y$ but $f(x) \neq f(y)$

$x \text{ --- } y$

$f(x) \text{ ~~~~~ } f(y)$

Term, $e ::= x | y | z | \dots | f(e) | g(e) | h(e) | \dots$

Equality constraints, $e_1 = e_2$

Ineq cts $e_1 \neq e_2$

Formula : $e_1 = e_2$ and $e_3 = e_4$ and \dots and $e_k = e_{k+1}$
and $e_{k+2} \neq e_{k+3}$ and $e_{k+4} \neq e_{k+5}$ and \dots

Congruence Closure

- ① Let the vertices $V =$ all subterms in φ
- ② For every equality constraint $e_1 = e_2$ in φ
draw an $e_1 - e_2$ edge
- ③ Repeat until no progress possible:
if I see $e_1 - e_2$ edge & $f(e_1)$ & $f(e_2) \in V$
then draw $f(e_1) - f(e_2)$ edge

④ Check if any inequality constraint $e_1 \neq e_2$ relates vertices in the same SCC.

⑤ sat if no, unsat if yes.

Ex Defining evenness (MSO with 1-successor)

$$\text{even}(x) = \exists y. x = y + y$$

$$\text{even}(x) = \forall E. \left(0 \in E \text{ and } \forall x. x \in E \Rightarrow x+2 \in E \right) \Rightarrow x \in E$$

- 0 is even

- x is even

$\Rightarrow x+2$ is even

- Smallest set with above properties

- DPLL(T)^{Theory}
 - Theory combination / Nelson-Oppen
-

Theory solver (T)

Input: Conjunction Σ -formula φ

Output: Yes \checkmark $\varphi \in T$, no otherwise

Ex: $x = y$ and $(y = z$ and $x \neq z)$ or $x = z$

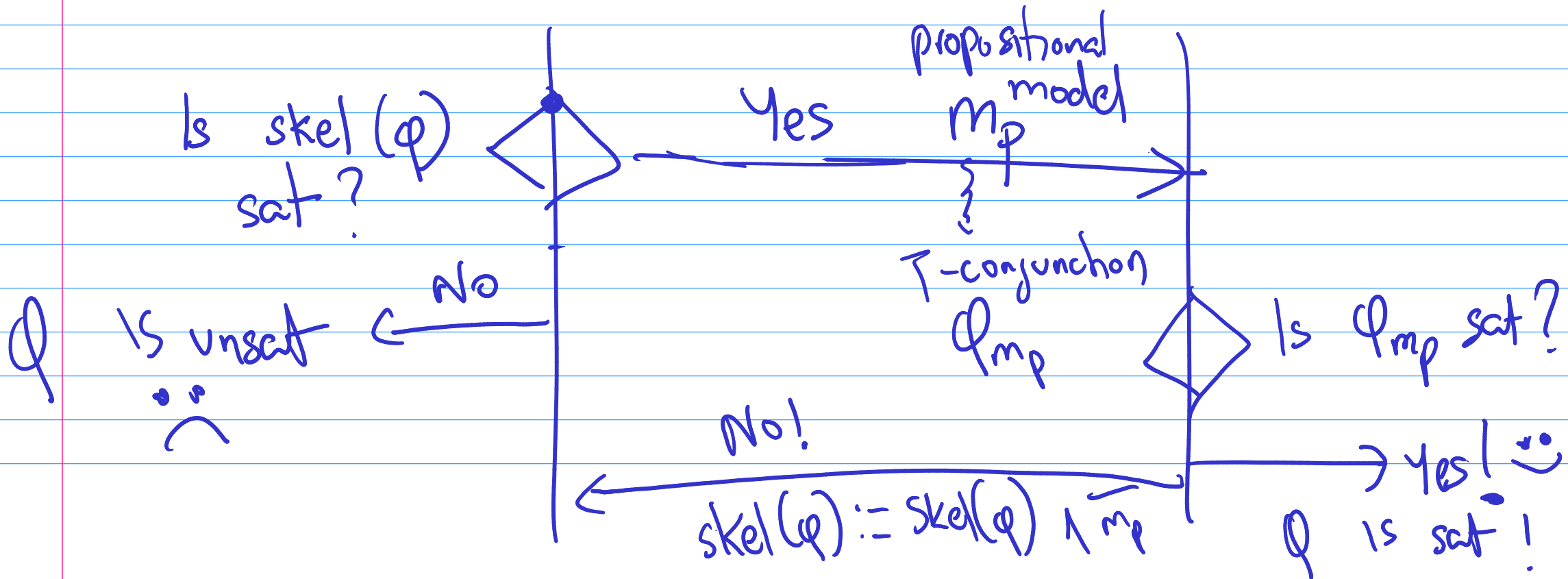
Propositional skeleton: v_1 and $(v_2$ and $v_3)$ or v_4

Propositional model: $v_1 \mapsto \text{true}$, $v_2 \mapsto \text{true}$, $v_3 \mapsto \text{true}$, $v_4 \mapsto \text{true}$
 $x = y$ and $y = z$ and $x \neq z$ and $x = z$

DPLL(T) Version 1.0

Sat solver

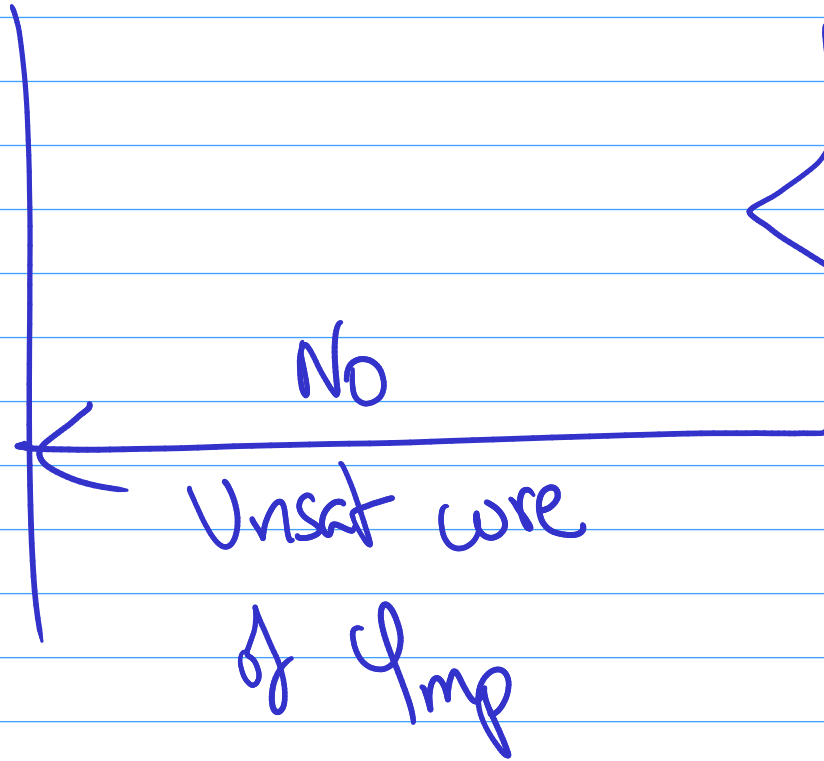
T-solver



DPLL(τ)

Version 2.0
Sat solver

T-solver



Is ϕ_{mp} sat?

DPLL(T) Version 3.0

Theory propagation: $q \Rightarrow x=y$

or $q \Rightarrow x \neq y$

Learn theory lemmas as the DPLL procedure proceeds.